Indexing, Storing and Retrieving Data in Neural Networks

Vincent Gripon







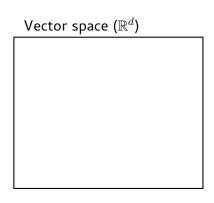
July 1st, 2017

Outline

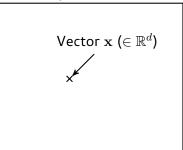
- Computer Vision and Neural Networks
- 2 Hopfield Neural Networks
- Willshaw Neural Networks
- 4 Conclusion

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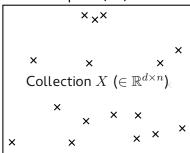
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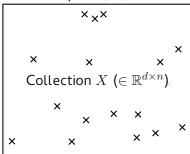
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- Unsupervised learning,
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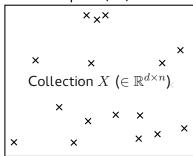
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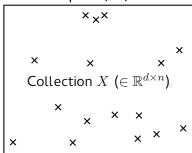
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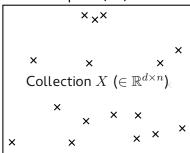
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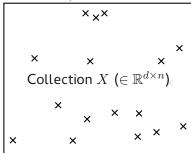
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```
0.88461 0.52899 0.39796 0.156
                              0.22615 0.16447 0.84366 0.7841 0.05846 0.3335 0.04146 0.79913
0.33711 0.22133 0.46954
                       0.8197
                              0.97514 0.79205 0.19736 0.33366 0.05208 0.04472 0.23042 0.93124
       0.45831 0.87231 0.71693 0.63678 0.54683 0.24892 0.32603 0.82655 0.08347 0.76076 0.59149
0.33481 0.62067 0.78087 0.55115 0.82832 0.46957 0.5429
                                                      0.7357 0.49622 0.09038 0.59702 0.38432
0.7534
       0.19463 0.41368 0.23335 0.01205 0.18668 0.9122
                                                      0.00722 0.64043 0.78145 0.94182 0.77094
0.45204 0.64851 0.0368 0.38763 0.99484 0.14494 0.76273 0.27692 0.33253 0.70724 0.7361 0.36882
0.35962
       0.0953 0.47678 0.92337 0.72545 0.3611
                                              0.05582 0.48013 0.5318 0.27792 0.90964 0.15971
 0.528
        0.4521
                0.6933
                       0.3117 0.57884 0.00188 0.06187 0.60576 0.94542 0.62769 0.82405 0.40215
0.4817
        0.3089 0.50847 0.56479 0.91013 0.38911 0.1955 0.19717 0.80548 0.0926 0.54935 0.2212
0.2007
       0.39793 0.76196 0.40977 0.5557 0.13638 0.11624 0.72516
                                                             0.711
                                                                     0.37856 0.34254 0.67796
0.18808
        0.495
               0.61931 0.85258 0.15338 0.95236 0.7579
                                                      0.83098 0.89072 0.30334 0.79318 0.93652
0.73792 0.10391 0.66104 0.11888 0.31796 0.11823 0.3503
                                                      0.21704 0.67531 0.10696 0.15614 0.88287
               0.7498 0.49826 0.56987 0.82922 0.50221 0.17014 0.14153 0.50203 0.71329 0.5883
0.87881
        0.5232
0.82059 0.15565 0.77045 0.65742 0.69325 0.81161 0.6689
                                                       0.2689
                                                              0.3157
                                                                     0.30891 0.10176 0.50745
0.35197 0.55392 0.10223 0.35126 0.41571 0.52171 0.06691 0.90193
                                                              0.5367
                                                                     0.32317 0.99646 0.32294
0.22276 0.84074 0.81502 0.19206 0.1831 0.95694 0.19064
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                                                              0.00731
                                                                      0.2918 0.75076 0.04846
0.83913 0.02509 0.23013 0.56805 0.48898 0.59413 0.00085 0.5568 0.94305 0.58304 0.97273 0.23122
0.8672 0.99681 0.55971 0.3113 0.93437 0.85008 0.31642 0.69461 0.92604 0.10722 0.9894 0.79062
0.73033 0.55927 0.65729 0.24986 0.83891 0.9922
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       0.26044 0.84335 0.92747 0.1045
                                       0.4189
                                              0.38507 0.61227 0.37956 0.96606 0.76146 0.66606
0.80537 0.69485
                0.216
                       0.07784 0.07887 0.0033
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0.58958 0.92455
               0.1638 0.98196 0.19221 0.06602 0.34438 0.62881 0.6117
                                                                      0.7972 0.42224 0.32251
0.91573 0.59866 0.83585 0.58174 0.16328 0.43048 0.93353 0.46086 0.84715 0.23059 0.77397 0.77893
               0.2249 0.59573 0.66231 0.51534
                                               0.836
                                                      0.75488 0.11812 0.52908 0.59819 0.46626
0.9289
       0.79767
0.62892 0.8288
               0.0395
                                                                             0.76498 0.70155
```













- Features extraction
- Features comparison

25% error rate in 2011





Learning

To learn is to **generalize** (\neq memorize),

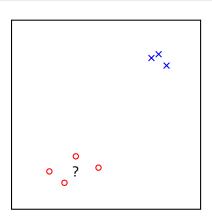
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- Requires an expert,
- Has a lot of applications:
 - Playing games,
 - Pattern recognition,
 - Attending to a lesson...



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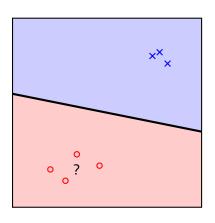
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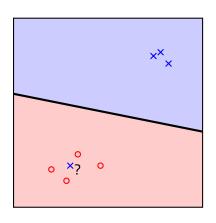
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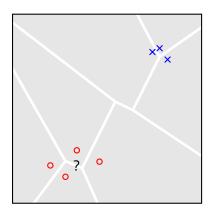
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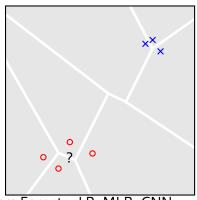


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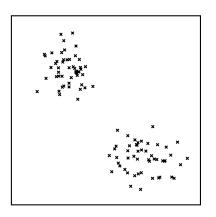
Supervised learning

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Classical methods: SVM, k-NN, Random Forests, LR, MLP, CNN...

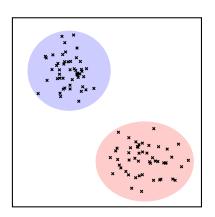
Unsupervised learning



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- Partitioning,
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- Many think this is the true support of intelligence:
 - Efficient representations, language,
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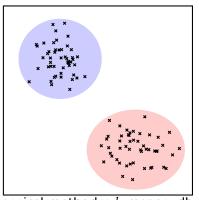
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Classical methods: k-means, db-scan, Kohonen maps, autoencoders, EM...

Definition

Given a collection $X \in \mathbb{R}^{d \times n}$ and a query vector \mathbf{x} :

- ① Is $x \in X$?
- ② Do we have $\mathbf{x}' \in X$ s.t. $\mathbf{x}' \approx \mathbf{x}$?

Exhaustive search

- Pros: no error, simple, concurrent,
- Cons: linear with both d and n.

- n = 1,000,000,000
- d = 128,
- 10,000 tests,
- On my laptop, takes approximatively 4 years.



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Hash tables (exact)

- Store items in an array of lists,
- Access array addresses using hash functions

Bloom filters (approximate)

- Retain already-seen hashes,
- Compare with probed ones.

- Use smooth hashes,
- Convert Euclidean to Hamming.

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Locality Sensitive Hashing (LSH)

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Search

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Given a collection $X \in \mathbb{R}^{d \times n}$ and a query vector \mathbf{x} , find:

$$\mathbf{x}' = \arg\min_{\mathbf{x}' \in X} \|\mathbf{x} - \mathbf{x}'\|,$$

given some metric.

Methods

- Exhaustive search again,
- Act on n and/or d:
 - ullet On n, partition the search space (problems with high dimensions),
 - On d, quantify the collection and/or the probe (e.g. Product Quantization).



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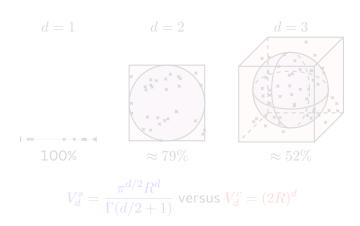
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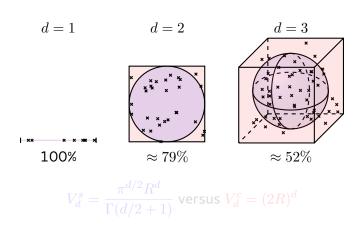
Curse of dimensionality

"Intuition is wrong in high dimension."



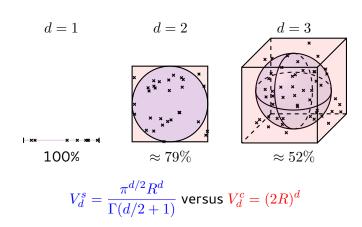
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A note on Neural Networks methods for vision

Definition

A neural network performs iteratively the concatenation of a linear and a nonlinear function.

$$\mathbf{y} = h_1(W_1(h_2(W_2 \dots h_i(W_i \mathbf{x})))).$$

Nonlinear functions

- Sigmoids (e.g. $x \mapsto 1/(1 + \exp(-x))$),
- Relus (e.g. $x \mapsto \max(0, x)$),
- Winner-Takes-All (WTA)...

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Framework

- $\mathbf{x} \in \{-1, 1\}^d$, $X \subset \{-1, 1\}^{d \times n}$,
- Example:

Retrieve it from -11-111-1?1

•
$$W = \sum_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top} - diag(\mathbf{x} \mathbf{x}^{\top}) = XX^{\top} - diag(XX^{\top}),$$
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- The update can be sequential (one coordinate at a time),
- $U(\mathbf{x}) \triangleq u(u(u(u(\dots u(\mathbf{x}))))).$

Framework

• $\mathbf{x} \in \{-1, 1\}^d$, $X \subset \{-1, 1\}^{d \times n}$,

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- Storing binary message
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Framework

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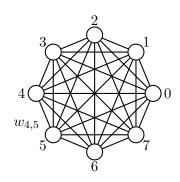
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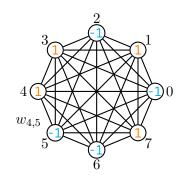
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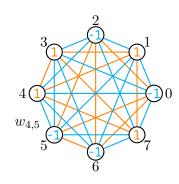
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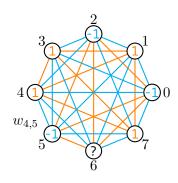
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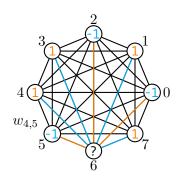
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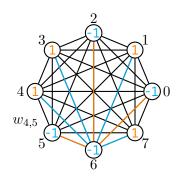
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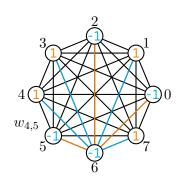
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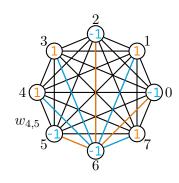
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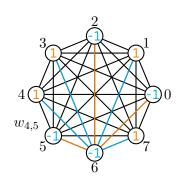
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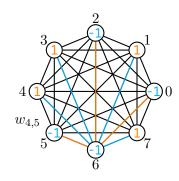
- $\mathbf{x} \in \{-1, 1\}^d$, $X \subset \{-1, 1\}^{d \times n}$,
- Example:
 - Storing binary message -11-111-1-11
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- $W = \sum_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top} diag(\mathbf{x} \mathbf{x}^{\top}) = XX^{\top} diag(XX^{\top}),$
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Theorem [1]

Consider $n = \frac{d}{\gamma \log(d)}$:

- If $\gamma > 6$, then for $d \to \infty$, $\mathbb{P}[\liminf_d \{ \cap_{\mathbf{x} \in X} \{ U(\mathbf{x}) = \mathbf{x} \}] = 1$
- If $\gamma > 4$, then $\mathbb{P}[\cap_{\mathbf{x}} \{U(\mathbf{x}) = \mathbf{x}\}] \to 1$.

- $\binom{d}{2}$ connections with n+1 possible values each \Rightarrow takes $\binom{d}{2} \log_2(n+1)$ bits without compression,
- To be compared to the entropy of $X \approx nd$ (Why \approx ?)
- When patterns are stable, we obtain $\eta \leq \frac{1}{2\log(d)\log_2(n+1)}$.
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Theorem [1]

Consider $\rho \in [0,1/2[$ and $n=(1-2\rho)^2\frac{d}{\gamma\log(d)}.$ $\tilde{\mathbf{x}}$ is such that it contains at most ρd symbols different from \mathbf{x} , then:

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Hamiltonian

The quantity $H(\mathbf{x}) = -\frac{1}{d} \sum_{i,j=1}^{d} W_{ij} \mathbf{x}_i \mathbf{x}_j$ is nonincreasing with u. More generally, $U(\mathbf{x})$ is a local minimum for H.



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Demonstration

- Create network,
- Test stability,
- Test retrievability.

Outline

- Computer Vision and Neural Networks
- Hopfield Neural Networks
- Willshaw Neural Networks
- 4 Conclusion

Willshaw Neural Networks

Framework

• $\mathbf{x} \in \{0, 1\}^d$, $\|\mathbf{x}\|_0 \ll d$, $X \subset \{0, 1\}^{d \times n}$.

- Example:
 - Storing binary message 01011001
 Patrious it from 01021021
 - Retrieve it from 010?10?1
- $W = \max_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top} \in \{0, 1\}^{d \times d}$ (= XX^{\top} for min-max-algebra)
- The update can be sequential (one coordinate at a time).
- $U(\mathbf{x}) \triangleq u(u(u(u(\dots u(\mathbf{x}))))).$

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 - 6

Willshaw Neural Networks

Framework

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Willshaw Neural Networks

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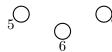
6

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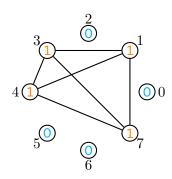




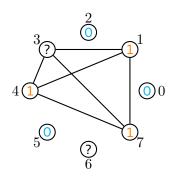




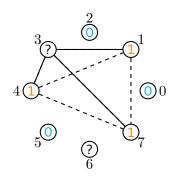
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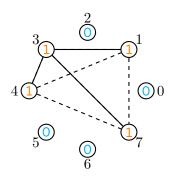
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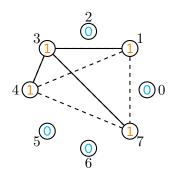
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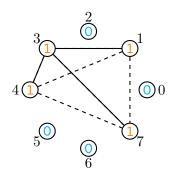
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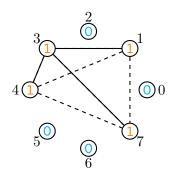
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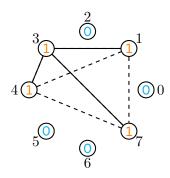
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Stability of stored vectors

Theorem [2]

Consider X generated with $\|\mathbf{x}\|_0 = \lfloor \log(d)/d \rfloor$ and \mathbf{x} chosen at random such that $\|\mathbf{x}\|_0 = \lfloor \log(d)/d \rfloor$. With $n = \alpha d^2 \log \log(d)/\log^2(d)$:

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Memory efficiency

Memory efficiency

- $\binom{d}{2}$ connections with 2 possible values each \Rightarrow takes $\binom{d}{2}$ bits without compression,
- To be compared to the entropy of *X*:

$$\approx ndH_2(\log(d)/d).$$

• When patterns are stable, we obtain

$$\eta \ge \frac{\alpha d(\log\log(d))^2}{2\log(d)} \to +\infty$$

Why?



Retrievability of stored vectors

Theorem [2]

Consider $n=\alpha d^2/\log^2(d)$, $\rho\in[0,1[$ such that $\lfloor\rho\log(d)\rfloor$ of 1s in ${\bf x}$ are erased to obtain $\tilde{{\bf x}}$. Then:

- If $\alpha < -\log(1-\exp(-1/(1-\rho)))$, then $\mathbb{P}[u(\tilde{\mathbf{x}}) = \mathbf{x}] \to 1$
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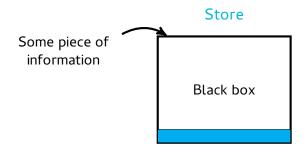
Demonstration

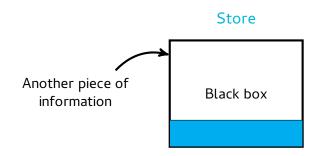
- Create network,
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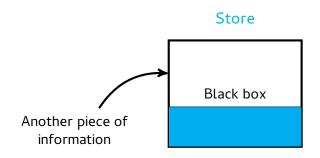
Outline

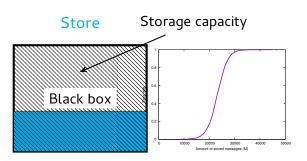
- Computer Vision and Neural Networks
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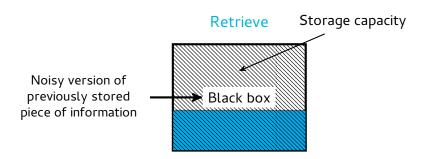


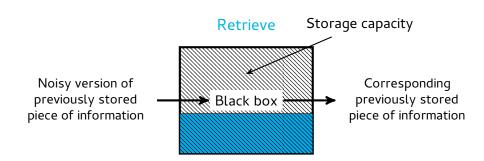


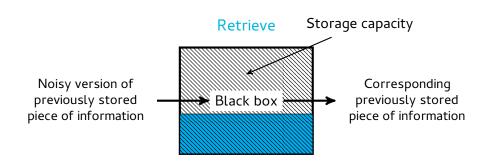












Piece of information = message, retrieve = decode associative memory = universal decoder

	Hopfield	Willshaw
Framework	$\mathbf{x} \in \{-1, 1\}^d$	$\mathbf{x} \in \{0,1\}^d, \ \mathbf{x}\ _0 \ll d$
Memory	$\mathbf{x}\mathbf{x}^\top - diag(\mathbf{x}\mathbf{x}^\top)$	$\mathbf{x}\mathbf{x}^{\top}$
Aggregation	$W = \sum_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top} - diag(\mathbf{x} \mathbf{x}^{\top})$	$W = \max_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top}$
Search	$u(W \cdot \tilde{\mathbf{x}})$	$u(W \otimes \tilde{\mathbf{x}})$

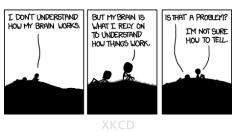
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	Hopfield	Willshaw
Framework	$\mathbf{x} \in \{-1, 1\}^d$	$\mathbf{x} \in \{0,1\}^d$, $\ \mathbf{x}\ _0 \ll c$
Memory	$\mathbf{x}\mathbf{x}^\top - diag(\mathbf{x}\mathbf{x}^\top)$	$\mathbf{x}\mathbf{x}^{\top}$
Aggregation	$W = \sum_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top} - diag(\mathbf{x} \mathbf{x}^{\top})$	$W = \max_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top}$
Search	$u(W \cdot \tilde{\mathbf{x}})$	$u(W \otimes \tilde{\mathbf{x}})$

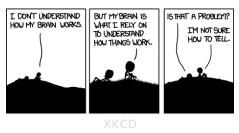
Framework $\mathbf{x} \in \{-1,1\}^d$ $\mathbf{x} \in \{0,1\}^d$, $\|\mathbf{x}\|_0 \ll d$ Memory $\mathbf{x}\mathbf{x}^\top - diag(\mathbf{x}\mathbf{x}^\top)$ $\mathbf{x}\mathbf{x}^\top$ Aggregation $W = \sum_{\mathbf{x} \in X} \mathbf{x}\mathbf{x}^\top - diag(\mathbf{x}\mathbf{x}^\top)$ $W = \max_{\mathbf{x} \in X} \mathbf{x}\mathbf{x}^\top$

	Hopfield	Willshaw
Framework	$\mathbf{x} \in \{-1, 1\}^d$	$\mathbf{x} \in \{0,1\}^d$, $\ \mathbf{x}\ _0 \ll d$
Memory	$\mathbf{x}\mathbf{x}^\top - diag(\mathbf{x}\mathbf{x}^\top)$	$\mathbf{x}\mathbf{x}^{\top}$
Aggregation	$W = \sum_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top} - diag(\mathbf{x} \mathbf{x}^{\top})$	$W = \max_{\mathbf{x} \in X} \mathbf{x} \mathbf{x}^{\top}$
Search	$u(W\cdot ilde{\mathbf{x}})$	$u(W \otimes \tilde{\mathbf{x}})$

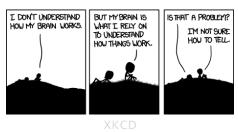
- Neural networks can do much more than learning,
- Neural networks are not just big mathematical functions
- Storing and indexing boils down to Gram matrices and strange algebras.



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